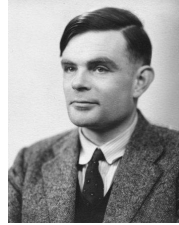
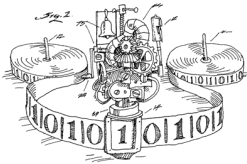


## Church-Turing Thesis



## Church-Turing Thesis



## Motivation: Encoding values with functions

```
t = lambda a: lambda b: a
f = lambda a: lambda b: b

def py_pred(p):
    return p(True)(False)

def f_not(p):
    """
    >>> py_pred(f_not(t))
    False
    >>> py_pred(f_not(f))
    True
    """
    return lambda a: lambda b: p(b)(a)
```

Exercise: Implement `f_and`, `f_or`  
`f_and = lambda p1: lambda p2:`  
`p2:` `p1(p2)(f)`  
`f_or = lambda p1: lambda p2:`  
`p2:` `p1(t)(p2)`

## A Cleaner Syntax

$\lambda[\text{argument}].[\text{return value}]$

$$f(x) \equiv fx$$

$$(\lambda x.M)N = M[x := N]$$

## Basic Functions

$$I = \lambda t.t$$

$$C_r = \lambda s.r$$

$$C_I = \lambda s.(\lambda t.t)$$

$$K = \lambda r.(\lambda s.r)$$

Exercise:  
 Confirm that `K I` is equivalent to `Cr`  
 Write down `Cr`, also write it in terms of `K` and `I`  
 Evaluate `(K K) (K K)`  
<http://www.chenyang.co/lambda>

## Currying

$$\begin{aligned}fgh &\equiv ((fg)h) \\ \lambda x.(\lambda y.xy) &\equiv \lambda xy.xy \\ \pi_1 &= \lambda xy.x \\ \pi_2 &= \lambda xy.y\end{aligned}$$

What does pi\_1 and pi\_2 remind you of?

## True and False

```
true = lambda a: lambda b: a
false = lambda a: lambda b: b
```

$T \equiv \lambda ab.a$

$F \equiv \lambda ab.b$

```
f_not = lambda p: lambda a:
        lambda b: p(b)(a)
f_and = lambda p1: lambda p2: p1(p2)(false)
f_or = lambda p1: lambda p2: p1(true)(p2)
```

Exercise:

Translate boolean operators into lambda notation

Not  $\equiv \lambda pab.pba$

And  $\equiv \lambda pq.pqF$

Or  $\equiv \lambda pq.pTq$

Write an if-function using lambda notation

## Zero, One, Two, ...

$$\begin{aligned}nf \rightarrow f^n &= f \circ f \circ \dots \circ f & 0 &\equiv \lambda fx.x \\ &= \lambda x.f(f(\dots f(x)\dots)) & 1 &\equiv \lambda fx.fx \\ n &= \lambda fx.f(f(\dots f(x)\dots)) & 2 &\equiv \lambda fx.f(fx) \\ Sn &= \lambda fx.f^{n+1}x \\ &= \lambda fx.f(f^n x) \\ &= \lambda fx.f(nfx) \\ S &= \lambda nfx.f(nfx)\end{aligned}$$

Exercise:

Implement add, mul, power

Suppose we have an operator pred that returns the predecessor of a number, implement subtraction using pred

Implement pred (Challenging)

## Recursion!

```
fac = lambda n: 1 if n == 0 else n * fac(n-1)
```

```
F = lambda fac: lambda n: 1 if n == 0 else n * fac(n-1)
```

$F(\text{fac}) = \text{fac}$  <- fac is the fixed point of F

$$F(YF) = YF$$

## Fixed Points

$$F(YF) = YF$$

$$\omega = (\lambda x.xx)(\lambda x.xx)$$

$$\begin{aligned}\omega' &= (\lambda x.F(xx))(\lambda x.F(xx)) \\ &= F((\lambda x.F(xx))(\lambda x.F(xx))) = F(\omega')\end{aligned}$$

$$Y = \lambda F.(\lambda x.F(xx))(\lambda x.F(xx))$$

## Puzzle

What is X? (There are 26 Ls)

Hint: A fixed point combinator has the form  $YF=F(YF)$

$$X = (\text{LLLLLLLLLLLLLLLLLLLLLLLLLLLLLLLL})$$

$L = \text{abcdefghijklmnopqrstuvwxyzr}$  (this is a fixed point combinator)

$$X = \lambda r.r(Xr)$$

$$XF = F(XF)$$

## **Conclusion + Acknowledgements**

---

Everything + more:

<http://xuanji.appspot.com/isicp/lambda.html>

Lambda interpreter:

<https://github.com/tarao/LambdaJS>

---